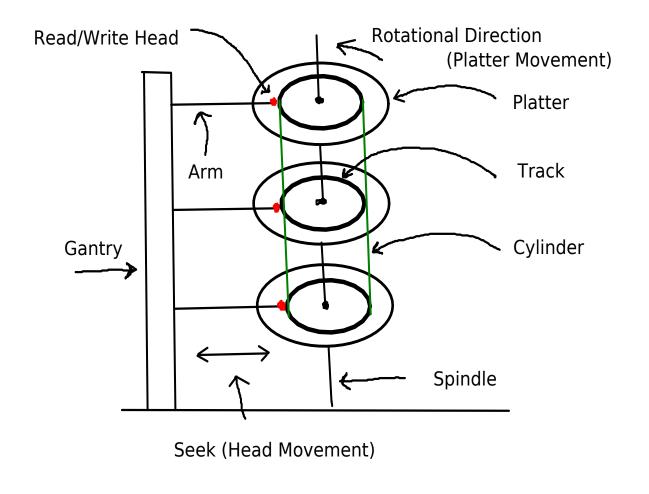
#### Computer Science CSCI 251

#### **Systems and Networks**

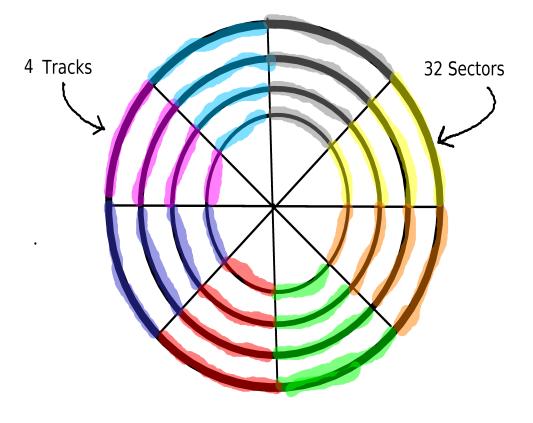
Dr. Peter Walsh
Department of Computer Science
Vancouver Island University
peter.walsh@viu.ca

### Disk Architecture



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# Disk Architecture cont.



### Terminology

- drives have 1 or more platters
- each platter has 1 or 2 sides
- each side has a read/write head
- each side is divided into concentric tracks
- platers and read/write heads are stacked
- a cylinder is the set of tracks the can be accessed simultaneously
- each track is divided into sectors
- moving a head over a track (or cylinder) is called seeking
- moving a head over a sector typically requires seeking and platter rotation

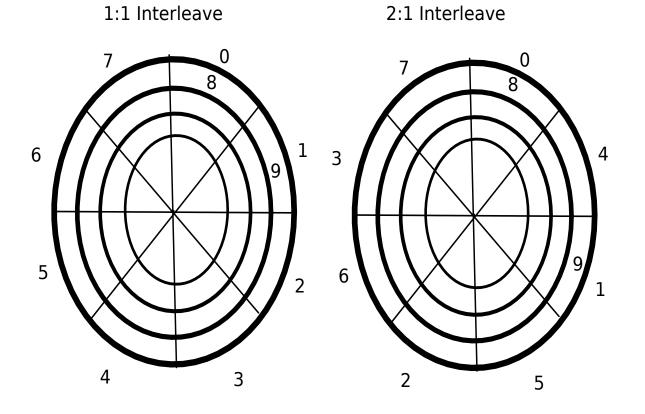
#### Sectors

- Size
  - smallest addressable part of the disk for read/write e.g., 512 bytes
  - Unix uses the term block (1 or more sectors)
- Address
  - cylinder#, head#, sector#
- Access Time
  - seek: move head to track/cylinder
  - rotational delay (latency): spin sector under the head
  - transfer: time to read/write data
- Old Drives
  - rotational delay (latency) << seek time</li>
- Modern Drives
  - ullet rotational delay (latency) pprox seek time

### Disk Architecture Optimizations

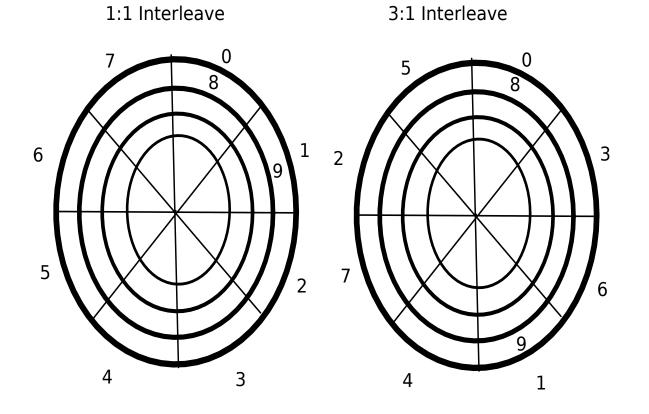
- Sector Interleaving
  - because of the overhead associated with read/write transfer, it may be more efficient to interleave sectors that are read in sequence
- Cylinder Skew
  - is a form of interleaving employed to decrease overhead when accessing consecutive sectors that cross track boundaries
- Multi Zoned
  - outer tracks are bigger and may be configured to contain more sectors that smaller inner tracks

### Sector Interleaving

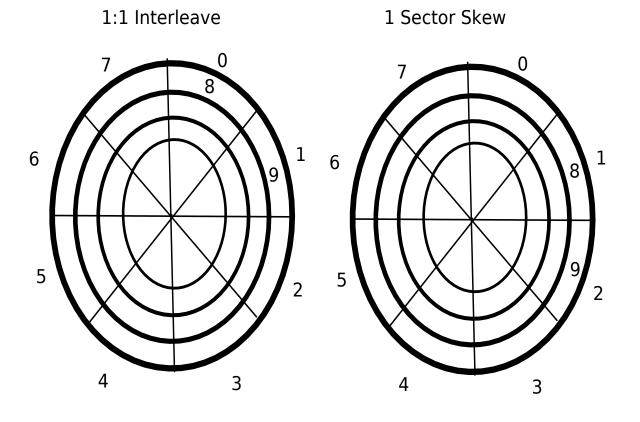


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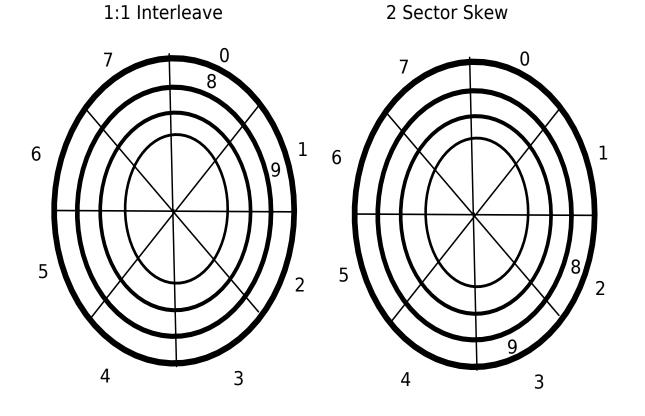
# Sector Interleaving cont.



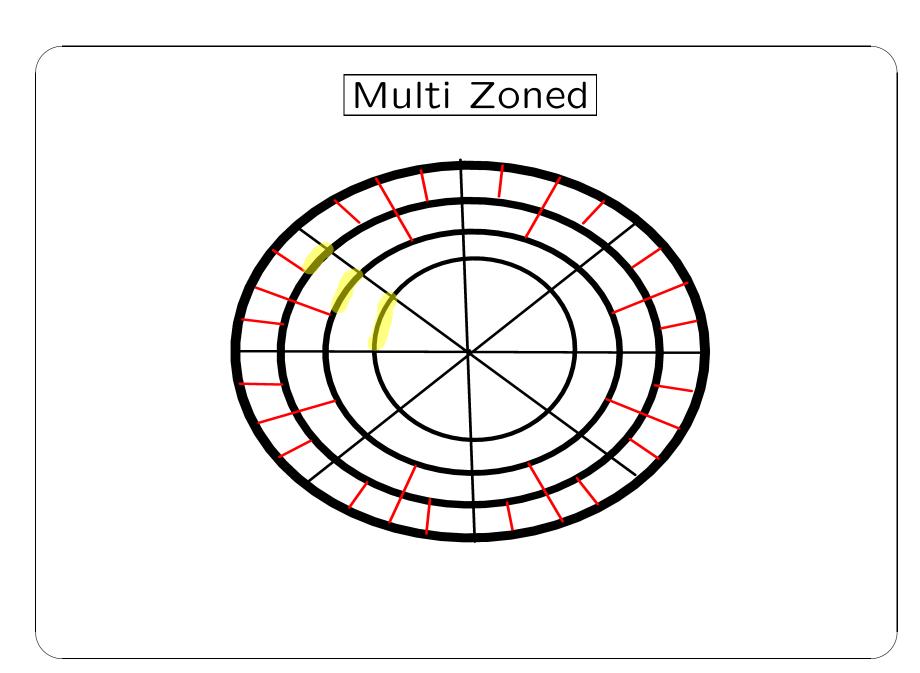
# Cylinder Skew



### Cylinder Skew cont.



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#### Workloads

- Random
  - issue reads to random locations on disk
  - e.g., database management system (DBMS)
- Sequential
  - read a large number of sectors consecutively
  - e.g., sequential file processing

# I/O Time

(Fig. 37.5 pp 463)	Cheetah	Barracuda	
RPM	15,000	7,200	
Average Seek	4ms	9ms	
Max Transfer	125 MB/s	105 MB/s	

Random Workload - 4 KB Reads Sequential Workload - 100 MB Read

### Random Workload (Cheetah)

## Random Workload (Cheetah)

## Random Workload (Barracuda)

Random Workload (Barracuda)

### Sequential Workload (Cheetah)

Tseek = 4ms

RPM = 15,000

1 rotation = 
$$\frac{60}{30,000}$$
 = 0.0025

=> Trotation = 2 ms

Max Rate = 125 MB/S

Ttransfer (100MB) =  $\frac{100}{125}$  s =  $\frac{600}{300}$  ms

## Sequential Workload (Cheetah)

Sequential Workload (Barracuda)

Tseek = 
$$9 \text{ ms}$$
  
RPM =  $7,200$   
 $\frac{1}{2}$  rotation =  $\frac{60}{14.400}$  =  $0.0041$ s  
 $\Rightarrow$  Trotation =  $4.1$  ms  
Max Rate =  $105$  MB/S  
Thansfer  $(100 \text{ MB})$  =  $\frac{100}{105}$  s =  $952$  ms

Sequential Workload (Barracuda)

$$T_{10} = T_{seek+1} + T_{transfer}$$

$$\approx 950 \text{ ms}$$

$$R_{10} = \frac{5}{T_{10}} = \frac{100}{950} \text{ MB/ms} = 105 \text{ mB/s}$$

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### Comparison Summary

(Fig. 37.6 pp 464)

Cheetah Barracuda

RPM 15,000 7,200

Average Seek 4ms 9ms

Max Transfer 125 MB/s 105MB/s

Random Workload - 4 KB Reads Sequential Workload - 100 MB Read

I/O Rate (Random) 0.66 MB/s 0.31 MB/s

I/O Rate (Sequential) 125 MB/s 105 MB/s

### Disk Scheduling

- Old Drives
  - rotational delay (latency) << seek time</li>
  - OS responsible for disk scheduling
- Modern Drives
  - ullet rotational delay (latency) pprox seek time
  - disk drive controller more involved in scheduling

### Disk Scheduling cont.

- Shortest Seek Time First (Scan)
  - sweep the disk servicing requests in-order across the tracks
- Shortest Position Time First
  - needs disk controller support
- OS/Controller Cooperation
  - OS selects a set of I/O requests to be serviced optimization is based on OS criteria
  - disk controller schedules requests for service optimization based on detailed knowledge of disk drive architecture and performance